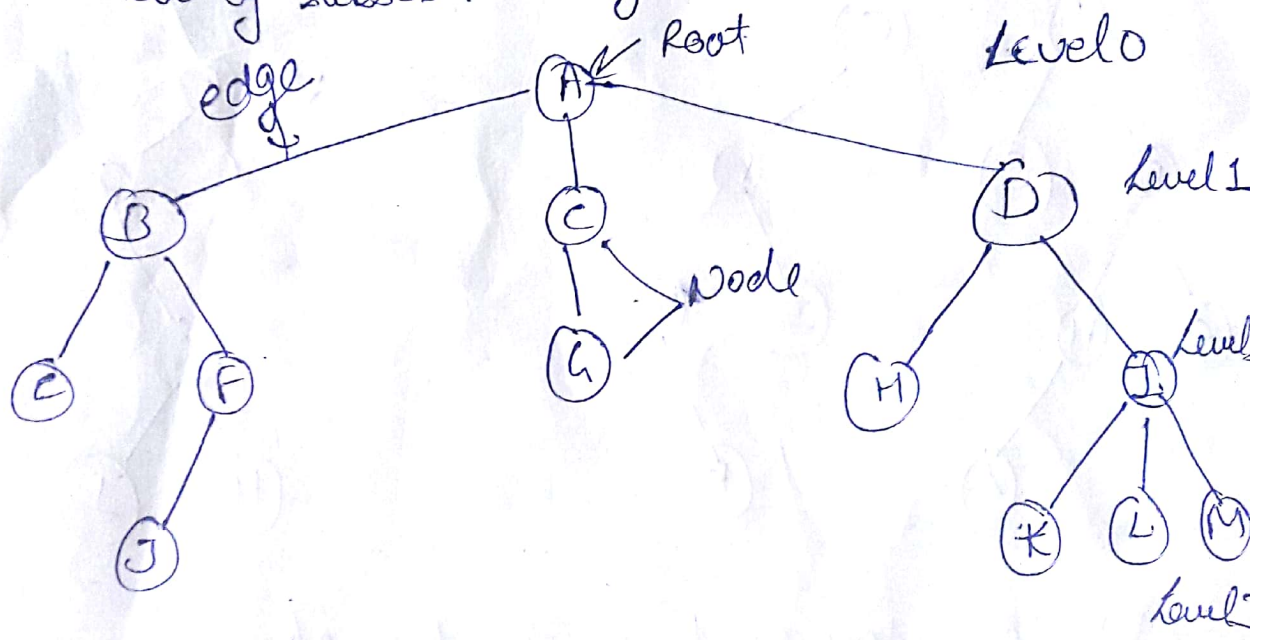


# Module-4

## TREE.

A tree is a non-linear data structure in which items are arranged in a sorted sequence. It is used to represent hierarchical relationship existing amongst several data items.

- There is a special data item called 'root' of tree.
- And remaining data items are partitioned into number of subsets, each of which is itself a tree.



### Terminology:

Root - First item of tree.

Node - Each data item.

Degree of Node - A - 3  
 B - 2  
 H - 0

Degree of Tree - Max. no. of nodes in a tree

Terminal/leaf Node - Node having no child.

Level - levels of child of tree.

Edge - connecting line of two nodes.

Depth - No. of Level + 1

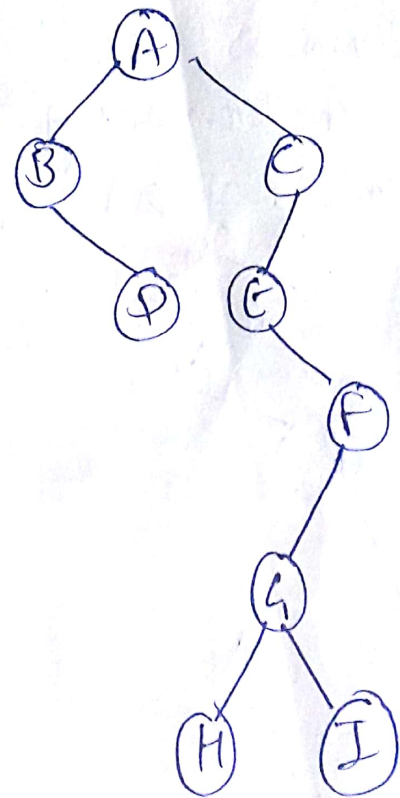
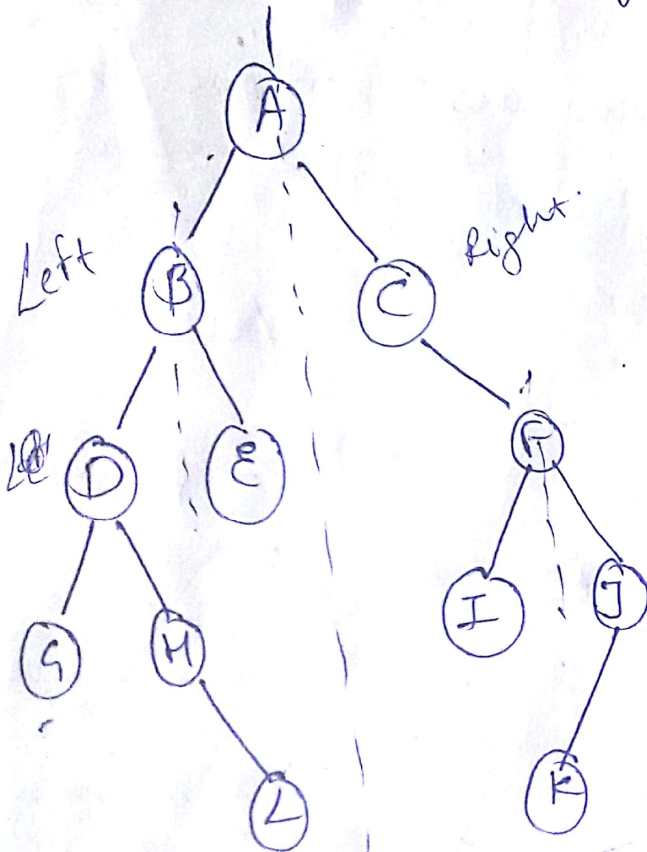
Forest - set of disjoint trees.

Traversing: shifting from one node to other.

→ Preorder - Root Left Right

→ Inorder - Left Root Right

→ Postorder - Left Right Root



→ Write orders of traversing for above trees.

Pre: A B D G H L C F I J K.

Post: G L H D E B I K J F C A.

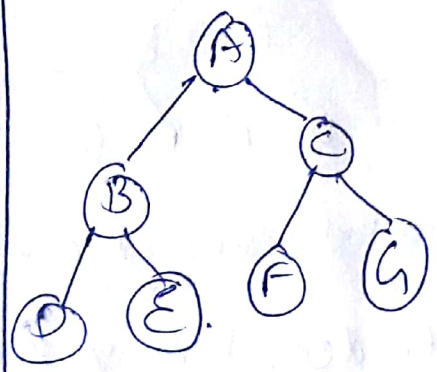
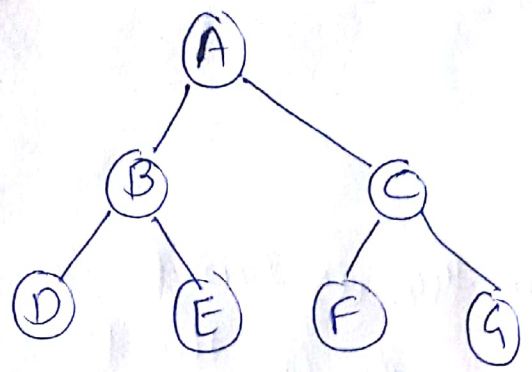
In: G D H L B E A C I F K J

create tree from following traversing results

(2)

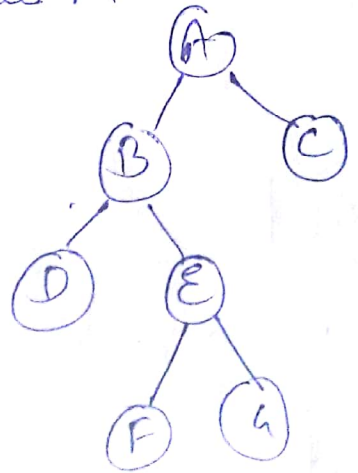
Pre - A B D E C F G  
 In - D B E A F C G

Post: DEB FGC A

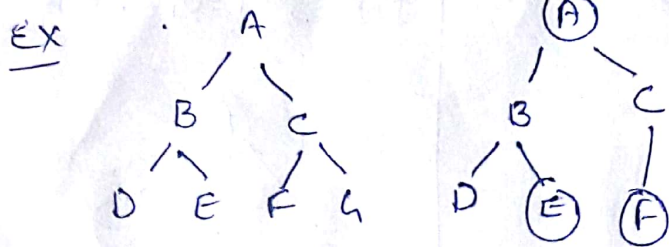


Binary Tree and Types

Strictly Binary Tree: If every non-terminal node in a binary tree consist of non-empty left subtree and right subtree, then such a tree is called strictly binary tree.

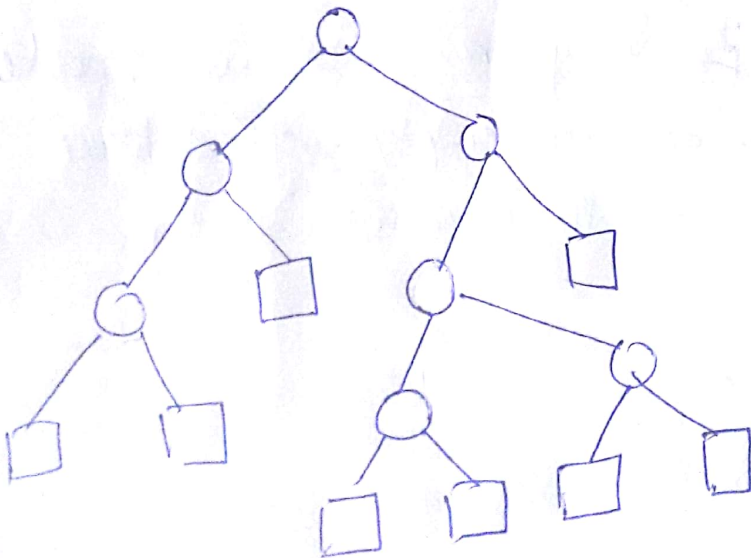


Complete Binary Tree - If all levels are completely filled except possibly the last level and last level has all keys as left as possible



Extended Binary Tree - A binary tree is said to be an extended binary tree if each node in the tree has either no child or exactly two children.

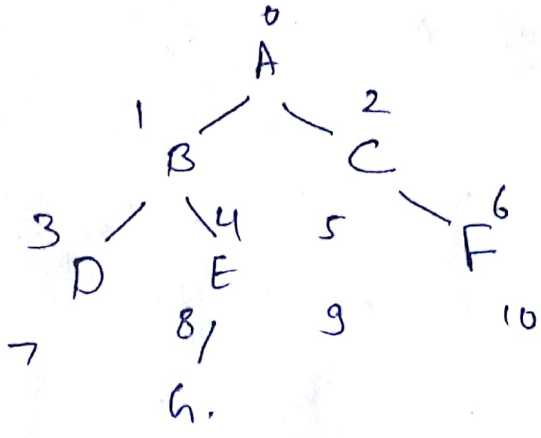
- Two children nodes - internal nodes ○
- no child → external nodes. □



Representation of trees in Memory:

- ① Array Representation
- ② Linked Representation.

# 1. Array Representation!



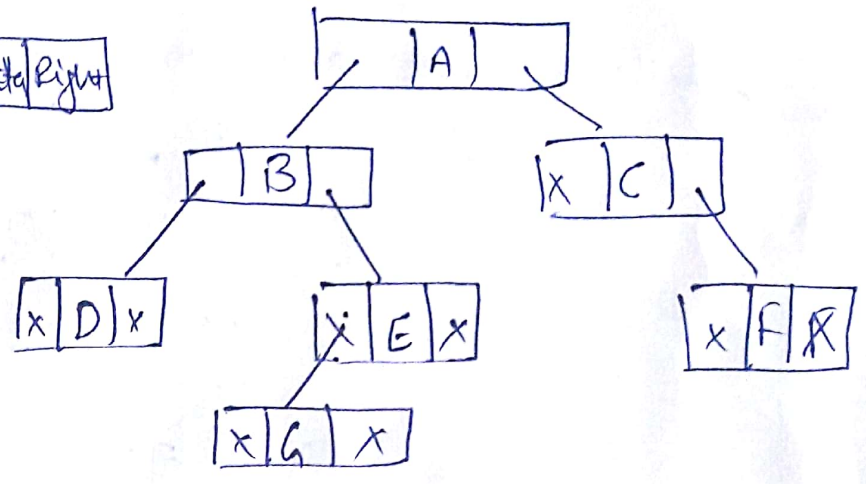
- 0 - A
- 1 - B
- 2 - C
- 3 - D
- 4 - E
- 5 - F
- 6 -
- 7 -
- 8 - G
- 9 -
- 10 -

# 2. Linked Representation.



```

struct node
{
  struct node * left;
  int data;
  struct node * right;
};
  
```



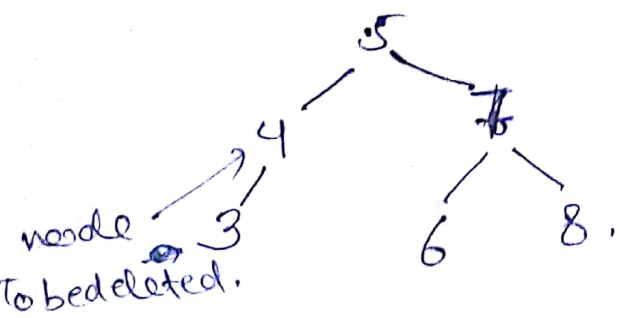
Binary Search Tree: It is a node based tree having

following properties:

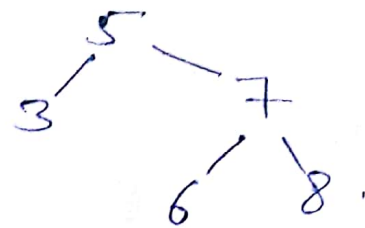
- > left subtrees have values less than root node
- right subtrees contains values greater than root node.
- If there are duplicate nodes, either avoid them or shift to right side of tree.

Ex- 8, 3, 6, 4, 7, 1, 10, 14, 13.

II. Having one child:

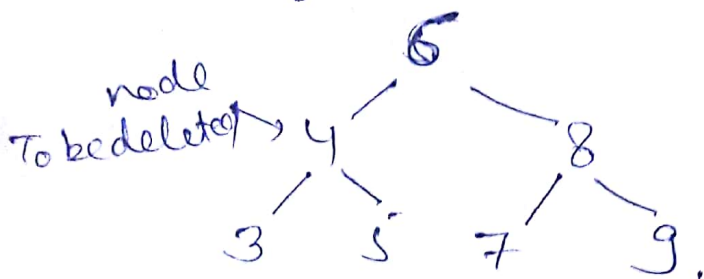


~~Swap~~ swap values of 3 with 4 & remove node 4



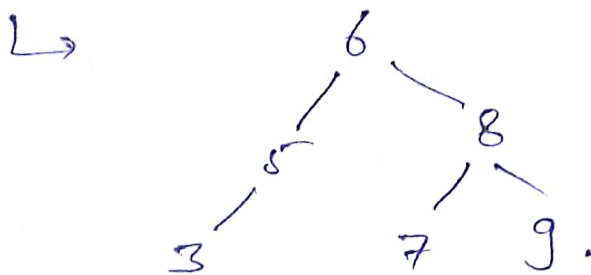
(4)

III. Having two children



I. find inorder traversal  
3 4 5 6 7 8 9.

II. next node after 4 is 5  
so, 5 will be root node.



# AVL Tree:

AVL tree is a self balanced tree. That means, an AVL tree is also a BST (binary search tree) but it is a balanced tree. It was introduced in the year of 1962 by G.M. Adelson-Velsky and E.M. Landis.

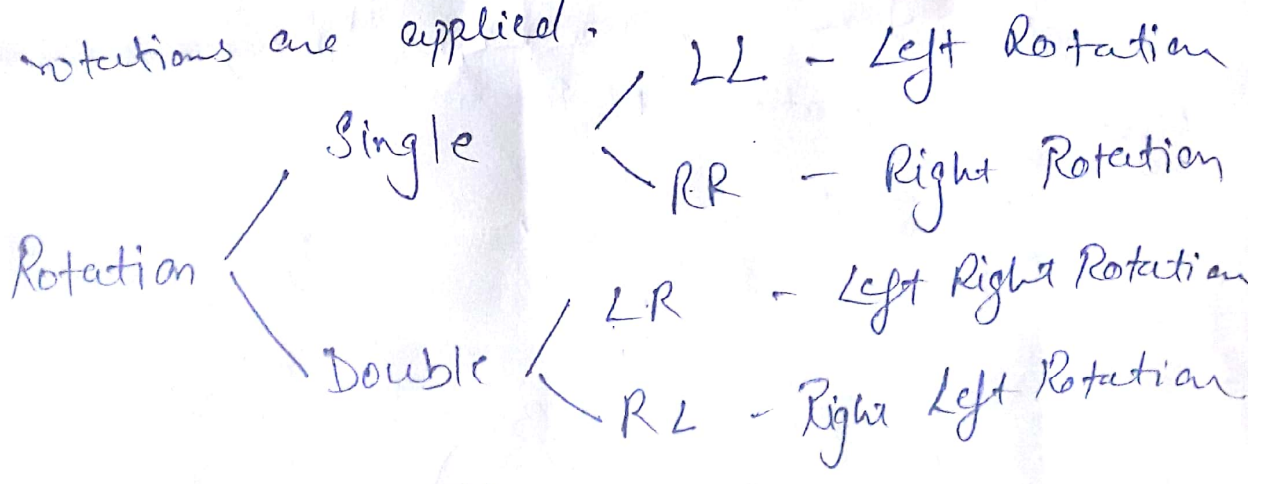
Every node having balance factor  $\boxed{+1, -1, \text{ or } 0}$

Balance factor = height of left subtree - height of right subtree

Every AVL is binary tree but not all binary tree are AVL.

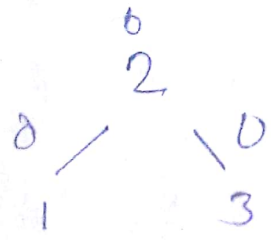
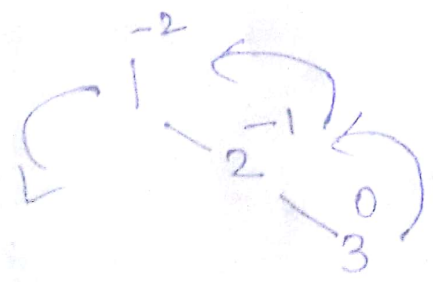
## AVL Tree Rotations:

In AVL tree performing every operation we need to check balance factor. If it is not balanced, then rotations are applied.



EX: 1, 2, 3

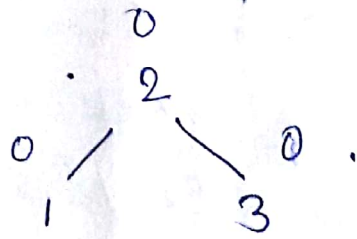
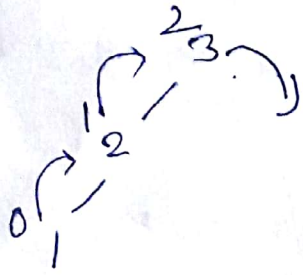
### LL Rotation



moves nodes to one position left.

## Single Right (RR) Rotation:

Instn 3, 2, 1.

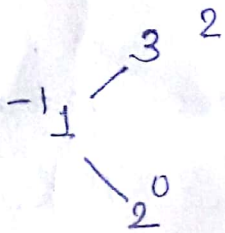


move node to one position right

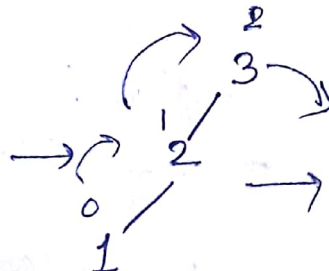
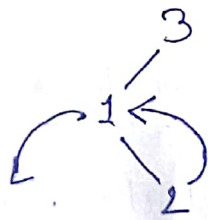
Balanced.

## Left Right (LR) Rotation:

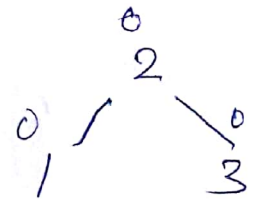
3, 1, 2



Applying LL Rotation



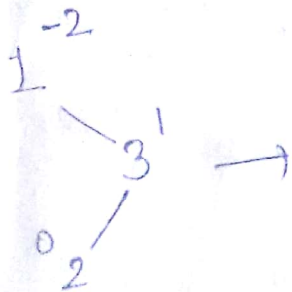
RR Rotation



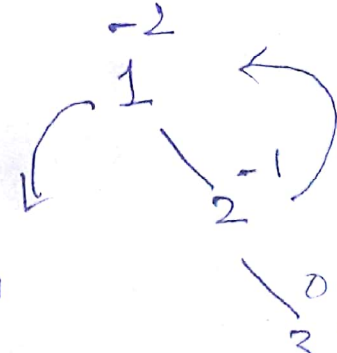
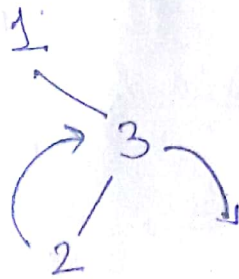
Balanced.

## Right Left Rotation: (RL Rotation)

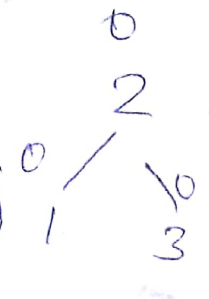
1, 3, 2



applying right rotation



apply Left rotation.



Balanced

→ Construct an AVL tree from numbers 1 to 8

Ex- 10, 20, 15, 25, 30, 16, 18, 19.

